Part 2: Design and implementation of an SCA core framework for a DSP platform

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Editor’s note: This is Part 2 of a two-part article. Part 1 ran in the March/April 2007 issue of Military Embedded Systems. You can read Part 1 online at: www.mil-embedded.com/articles/id/?2065.

The authors present the design and implementation of the SCA 2.2 Core Framework for a TI DSP platform and provide the rationale behind design decisions and initial profiling results.

Results

Profiling was performed on the framework and applications using two different metrics: memory footprint and cycle count. All results were obtained from a single-chip configuration. That is, all framework and waveform components were collocated within the same processor. Hence, these results do not include the effects of a transport layer. No optimization was performed in either the framework or the waveform components and they include debug information. It is very important to emphasize that these results represent initial measurements and are subject to further investigation, validation, and optimization.

Memory footprint

The total memory used by the system is little more than 1.46 MB, which represents less than 2 percent of the available memory per DSP (128 MB) in the platform. Table 1 shows the memory breakdown by major software components. The .ERAM$heap field represents the total memory available to serve dynamic memory allocation requests. The footprint contribution from support libraries (Generic Runtime Library, Math Library, and so on) is considered under the “Other” category. Figure 4 shows a graphical representation of the main components’ contribution to total memory allocation.

<table>
<thead>
<tr>
<th>SW Component</th>
<th>Footprint (Bytes)</th>
</tr>
</thead>
<tbody>
<tr>
<td>CF</td>
<td>556555</td>
</tr>
<tr>
<td>Parsers</td>
<td>31511</td>
</tr>
<tr>
<td>ORB</td>
<td>212412</td>
</tr>
<tr>
<td>Application</td>
<td>385624</td>
</tr>
<tr>
<td>.ERAM$heap</td>
<td>131072</td>
</tr>
<tr>
<td>Other</td>
<td>144067</td>
</tr>
<tr>
<td>Total memory</td>
<td>1461241</td>
</tr>
</tbody>
</table>

Table 1

In analyzing the Core Framework’s (CF’s) memory requirements, we found that almost 70 percent of total memory allocated for the CF came from the C++ mapping of the SCA CF IDL interfaces. It is important to note that the CF IDL descriptions contained all the interfaces defined in the SCA CF, including some that were not used in single-processor operation (for example, DeviceManager, Device). It is possible to optimize the C++ bindings of IDL interfaces by adding more control to the IDL compiler, enabling more selective code generation (such as for specific interfaces generate stub only, or skeletons only, or nothing). This approach opens the door for potentially large improvements, depending on how much of the IDL interfaces are used. This is a well-understood approach, although it was not implemented in this project. Another important qualifier for these results is the absence of Device-related interfaces. No DeviceManager or Device interfaces were implemented. The methods in DomainManager relative to Device and service registration were not implemented either.

The memory requirement results for the application include BPSK and QPSK components, along with Assembly Controller, Channel, Demodulator, Resource-
Factory, and the user interface. The main waveform components have very similar footprints as expected. However, the functionality of these components is extremely simple. More complex waveforms will require more memory.

**Performance profile**

CPU cycle requirements were collected from the CF’s startup tasks: domain initialization and waveform creation. The results are shown in Table 2. Domain initialization includes the instantiation of Domain Manager, ApplicationFactory, and ResourceFactory, which durations are independent of the waveform deployed. Waveform creation represents the execution of ApplicationFactory’s create(). It includes descriptor parsing, task scheduling and initialization, and component connection. Keep in mind that waveform creation is waveform-specific, and these results only apply to our test waveforms.

Opposite to initialization tasks, ORB performance had a great impact on the system throughput because all intercomponent communications were established using CORBA messages. Two specific scenarios were profiled:

- **Invocation**: Round-trip cycle count for a simple method invocation with no arguments
- **Marshaling**: Round-trip cycle count for a simple method invocation with basic arguments

Two different argument types were evaluated:

- Single data type
- Sequence (1,024 elements)

In our version of e*ORB, even for interfaces with no arguments in their IDL definitions, a CORBA::Environment variable must be sent as an argument because of the lack of exception support.

In both scenarios, client and server were launched as separate DSP/BIOS tasks with priorities 2 and 1, respectively. The e*ORB profiling results for different primitive data types are summarized in Tables 3 and 4.

![Figure 5](image_url)  
**Figure 5**  

It takes 4,208 clock cycles to make a round-trip marshalling call with a single float, while it takes 6,908 clock cycles to send a sequence with 1,024 floats. Averaging, it only takes 6.74 clock cycles to transfer each element in the sequence.

**Impact on data rate performance**

The framework overhead incurred during instantiation and waveform deployment can be arranged to happen off-line. The only aspect of the SCA that impacts system throughput is the dependency on CORBA for intercomponent communications. The maximum system data rate depends on many factors: algorithm processing delays, framework delays, analog-to-digital conversion rate, and so on. To isolate the impact of the framework, we used the results shown in Table 4 to estimate an upper bound for the system data rate.

Ignoring processing delays, the maximum achievable data rate is given by:

$$R_{max} = \frac{1}{T_m + T_{tr}}$$
where \( T_m \) is the delay due to middleware processing and \( T_{tr} \) is the delay due to transport mechanisms. In our system, we only considered \( T_{tr} \) because no transport mechanisms had been developed at the time.

The average delay per bit due to middleware message passing \( T_m \) is given by:

\[
T_m = \frac{D_t \cdot S_s}{N_p \cdot n}
\]

where \( D_t \) is the measured delay as shown in Table 4. \( S_s \) is the number of samples per symbol. \( N_p \) is the packet size, and \( n \) is the number of bits per symbol. To estimate the maximum data rate allowed by the framework, we assumed \( S_s = 8 \) and \( n = 1 \). The clock speed in our system is 720 MHz. Substituting these values into the expression for \( T_m \) for a single float type transfer that according to Table 4 takes 4,142 cycles, the maximum data rate achievable is \( R_{\text{max}} = 21,728 \) bits per second. However, if we consider sending a sequence of 1,024 floats, the transfer takes 6,072 clock cycles allowing a maximum data rate \( R_{\text{max}} = 15,177,865 \) bits per second. These results highlight the need for block processing within the SCA, trading off latency and performance.

The future of SCA on DSP

One of the main concerns of following the SCA is the heavy infrastructure required to support it. To ease requirements in terms of performance, cost, and power consumption, we propose an implementation of the SCA Core Framework for a TI C64 DSP platform. This approach is feasible thanks to the latest developments in software tools and ORB technology. By having an SCA core framework in a DSP, all the benefits in software reuse and deployment flexibility brought by the SCA can be achieved in a more efficient platform. In terms of performance, our complete implementation requires about 1.5 MB, which represents about 1 percent of the total memory available per DSP in our platform. Even though CORBA introduces some delays and overhead, the overall effect can be reduced by sending data packets instead of single elements. The source code for the framework and sample waveforms is available at http://ossie.mprg.org.

Finally, it also needs to be noted that SDR technology continues to rapidly evolve and improve. The work referenced herein was conducted and profiled about a year ago. If the design and implementation were repeated today, the results (such as performance, memory footprint, and so on) would be that much more impressive in validating the SCA on DSP.

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